

Vapnik-Chervonenkis Density

Matthias Aschenbrenner

University of California, Los Angeles

(joint with A. Dolich, D. Haskell, D. Macpherson, and S. Starchenko)

UCLA

VC dimension and VC density

Let (X, \mathcal{S}) be a set system, i.e., X is a set (the **base set**), and \mathcal{S} is a collection of subsets of X . (We sometimes also speak of a **set system** \mathcal{S} on X .)

Let (X, \mathcal{S}) be a set system, i.e., X is a set (the **base set**), and \mathcal{S} is a collection of subsets of X . (We sometimes also speak of a **set system** \mathcal{S} on X .)

Given $A \subseteq X$, we let

$$\mathcal{S} \cap A := \{S \cap A : S \in \mathcal{S}\}$$

and call $(A, \mathcal{S} \cap A)$ **the set system on A induced by \mathcal{S}** .

Let (X, \mathcal{S}) be a set system, i.e., X is a set (the **base set**), and \mathcal{S} is a collection of subsets of X . (We sometimes also speak of a **set system \mathcal{S} on X** .)

Given $A \subseteq X$, we let

$$\mathcal{S} \cap A := \{S \cap A : S \in \mathcal{S}\}$$

and call $(A, \mathcal{S} \cap A)$ **the set system on A induced by \mathcal{S}** .

We say A is **shattered by \mathcal{S}** if $\mathcal{S} \cap A = 2^A$.

If $\mathcal{S} \neq \emptyset$, then we define the **VC dimension of \mathcal{S}** , denoted by $\text{VC}(\mathcal{S})$, as the supremum (in $\mathbb{N} \cup \{\infty\}$) of the sizes of all finite subsets of X shattered by \mathcal{S} . We also decree $\text{VC}(\emptyset) := -\infty$.

Examples

Examples

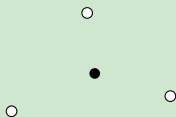
- 1 $X = \mathbb{R}$, $\mathcal{S} =$ all unbounded intervals. Then $VC(\mathcal{S}) = 2$.

Examples

- 1 $X = \mathbb{R}$, $\mathcal{S} =$ all unbounded intervals. Then $VC(\mathcal{S}) = 2$.
- 2 $X = \mathbb{R}^2$, $\mathcal{S} =$ all halfspaces. Then $VC(\mathcal{S}) = 3$.

Examples

- 1 $X = \mathbb{R}$, $\mathcal{S} =$ all unbounded intervals. Then $VC(\mathcal{S}) = 2$.
- 2 $X = \mathbb{R}^2$, $\mathcal{S} =$ all halfspaces. Then $VC(\mathcal{S}) = 3$.



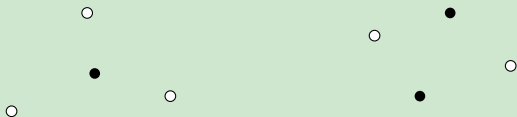
One point in the convex hull
of the others



No point in the convex hull
of the others

Examples

- 1 $X = \mathbb{R}$, $\mathcal{S} =$ all unbounded intervals. Then $VC(\mathcal{S}) = 2$.
- 2 $X = \mathbb{R}^2$, $\mathcal{S} =$ all halfspaces. Then $VC(\mathcal{S}) = 3$.



One point in the convex hull
of the others

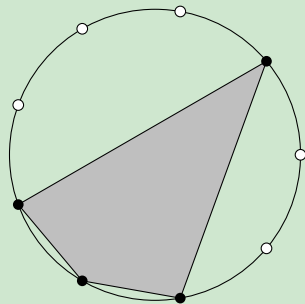
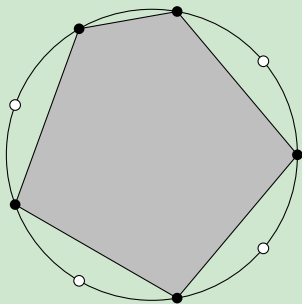
No point in the convex hull
of the others

- 3 Let $\mathcal{S} =$ half spaces in \mathbb{R}^d . Then $VC(\mathcal{S}) = d + 1$.
(The inequality \leq follows from *Radon's Lemma*.)

Examples (continued)

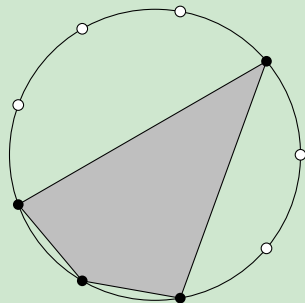
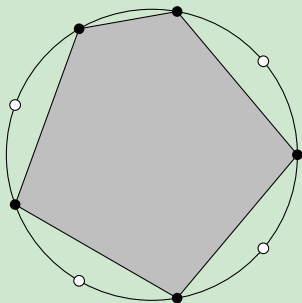
Examples (continued)

- 4 $X = \mathbb{R}^2$, $\mathcal{S} = \text{all convex polygons}$. Then $\text{VC}(\mathcal{S}) = \infty$.



Examples (continued)

- 4 $X = \mathbb{R}^2$, $\mathcal{S} = \text{all convex polygons}$. Then $\text{VC}(\mathcal{S}) = \infty$.



(But $\text{VC}(\{\text{convex } n\text{-gons in } \mathbb{R}^2\}) = 2n + 1$.)

The function

$$n \mapsto \pi_{\mathcal{S}}(n) := \max \left\{ |\mathcal{S} \cap A| : A \in \binom{X}{n} \right\} : \mathbb{N} \rightarrow \mathbb{N}$$

is called the **shatter function of \mathcal{S}** .

The function

$$n \mapsto \pi_{\mathcal{S}}(n) := \max \left\{ |\mathcal{S} \cap A| : A \in \binom{X}{n} \right\} : \mathbb{N} \rightarrow \mathbb{N}$$

is called the **shatter function of \mathcal{S}** . Then

$$\text{VC}(\mathcal{S}) = \sup \{n : \pi_{\mathcal{S}}(n) = 2^n\}.$$

One says that \mathcal{S} is a **VC class** if $\text{VC}(\mathcal{S}) < \infty$.

The function

$$n \mapsto \pi_{\mathcal{S}}(n) := \max \left\{ |\mathcal{S} \cap A| : A \in \binom{X}{n} \right\} : \mathbb{N} \rightarrow \mathbb{N}$$

is called the **shatter function of \mathcal{S}** . Then

$$\text{VC}(\mathcal{S}) = \sup \{n : \pi_{\mathcal{S}}(n) = 2^n\}.$$

One says that \mathcal{S} is a **VC class** if $\text{VC}(\mathcal{S}) < \infty$.

The notion of VC dimension was introduced by Vladimir Vapnik and Alexey Chervonenkis in the early 1970s, in the context of computational learning theory.



VC dimension and VC density

A surprising dichotomy holds for π_S :

A surprising dichotomy holds for $\pi_{\mathcal{S}}$:

The Sauer-Shelah dichotomy

Either

- $\pi_{\mathcal{S}}(n) = 2^n$ for every n (if \mathcal{S} is not a VC class),

or

- $\pi_{\mathcal{S}}(n) \leq \binom{n}{\leq d} := \binom{n}{0} + \dots + \binom{n}{d}$ where $d = \text{VC}(\mathcal{S}) < \infty$.

A surprising dichotomy holds for $\pi_{\mathcal{S}}$:

The Sauer-Shelah dichotomy

Either

- $\pi_{\mathcal{S}}(n) = 2^n$ for every n (if \mathcal{S} is not a VC class),

or

- $\pi_{\mathcal{S}}(n) \leq \binom{n}{\leq d} := \binom{n}{0} + \dots + \binom{n}{d}$ where $d = \text{VC}(\mathcal{S}) < \infty$.

One may now define the **VC density** of \mathcal{S} as

$$\text{vc}(\mathcal{S}) = \begin{cases} \inf\{r \in \mathbb{R}^{>0} : \pi_{\mathcal{S}}(n) = O(n^r)\} & \text{if } \text{VC}(\mathcal{S}) < \infty \\ \infty & \text{otherwise.} \end{cases}$$

A surprising dichotomy holds for $\pi_{\mathcal{S}}$:

The Sauer-Shelah dichotomy

Either

- $\pi_{\mathcal{S}}(n) = 2^n$ for every n (if \mathcal{S} is not a VC class),

or

- $\pi_{\mathcal{S}}(n) \leq \binom{n}{\leq d} := \binom{n}{0} + \dots + \binom{n}{d}$ where $d = \text{VC}(\mathcal{S}) < \infty$.

One may now define the **VC density** of \mathcal{S} as

$$\text{vc}(\mathcal{S}) = \begin{cases} \inf\{r \in \mathbb{R}^{>0} : \pi_{\mathcal{S}}(n) = O(n^r)\} & \text{if } \text{VC}(\mathcal{S}) < \infty \\ \infty & \text{otherwise.} \end{cases}$$

We also define $\text{vc}(\emptyset) := -\infty$.

Examples

① $\mathcal{S} = \binom{X}{\leq d}$. Then $\text{VC}(\mathcal{S}) = \text{vc}(\mathcal{S}) = d$; in fact $\pi_{\mathcal{S}}(n) = \binom{n}{\leq d}$.

Examples

- 1 $\mathcal{S} = \binom{X}{\leq d}$. Then $\text{VC}(\mathcal{S}) = \text{vc}(\mathcal{S}) = d$; in fact $\pi_{\mathcal{S}}(n) = \binom{n}{\leq d}$.
- 2 $\mathcal{S} = \text{half spaces in } \mathbb{R}^d$. Then $\text{VC}(\mathcal{S}) = d + 1$, $\text{vc}(\mathcal{S}) = d$.

Examples

- 1 $\mathcal{S} = \binom{X}{\leq d}$. Then $\text{VC}(\mathcal{S}) = \text{vc}(\mathcal{S}) = d$; in fact $\pi_{\mathcal{S}}(n) = \binom{n}{\leq d}$.
- 2 $\mathcal{S} =$ half spaces in \mathbb{R}^d . Then $\text{VC}(\mathcal{S}) = d + 1$, $\text{vc}(\mathcal{S}) = d$.

VC density is often the right measure for the combinatorial complexity of a set system. (E.g., it is related to packing numbers and entropy).

Examples

- 1 $\mathcal{S} = \binom{X}{\leq d}$. Then $\text{VC}(\mathcal{S}) = \text{vc}(\mathcal{S}) = d$; in fact $\pi_{\mathcal{S}}(n) = \binom{n}{\leq d}$.
- 2 $\mathcal{S} =$ half spaces in \mathbb{R}^d . Then $\text{VC}(\mathcal{S}) = d + 1$, $\text{vc}(\mathcal{S}) = d$.

VC density is often the right measure for the combinatorial complexity of a set system. (E.g., it is related to packing numbers and entropy).

Some basic properties:

- $\text{vc}(\mathcal{S}) \leq \text{VC}(\mathcal{S})$, and if one is finite then so is the other;

Examples

- 1 $\mathcal{S} = \binom{X}{\leq d}$. Then $\text{VC}(\mathcal{S}) = \text{vc}(\mathcal{S}) = d$; in fact $\pi_{\mathcal{S}}(n) = \binom{n}{\leq d}$.
- 2 $\mathcal{S} =$ half spaces in \mathbb{R}^d . Then $\text{VC}(\mathcal{S}) = d + 1$, $\text{vc}(\mathcal{S}) = d$.

VC density is often the right measure for the combinatorial complexity of a set system. (E.g., it is related to packing numbers and entropy).

Some basic properties:

- $\text{vc}(\mathcal{S}) \leq \text{VC}(\mathcal{S})$, and if one is finite then so is the other;
- $\text{VC}(\mathcal{S}) = 0 \iff |\mathcal{S}| = 1$;

Examples

- 1 $\mathcal{S} = \binom{X}{\leq d}$. Then $\text{VC}(\mathcal{S}) = \text{vc}(\mathcal{S}) = d$; in fact $\pi_{\mathcal{S}}(n) = \binom{n}{\leq d}$.
- 2 $\mathcal{S} = \text{half spaces in } \mathbb{R}^d$. Then $\text{VC}(\mathcal{S}) = d + 1$, $\text{vc}(\mathcal{S}) = d$.

VC density is often the right measure for the combinatorial complexity of a set system. (E.g., it is related to packing numbers and entropy).

Some basic properties:

- $\text{vc}(\mathcal{S}) \leq \text{VC}(\mathcal{S})$, and if one is finite then so is the other;
- $\text{VC}(\mathcal{S}) = 0 \iff |\mathcal{S}| = 1$;
- \mathcal{S} is finite $\iff \text{vc}(\mathcal{S}) = 0 \iff \text{vc}(\mathcal{S}) < 1$;

Examples

- 1 $\mathcal{S} = \left(\binom{X}{\leq d} \right)$. Then $\text{VC}(\mathcal{S}) = \text{vc}(\mathcal{S}) = d$; in fact $\pi_{\mathcal{S}}(n) = \binom{n}{\leq d}$.
- 2 $\mathcal{S} = \text{half spaces in } \mathbb{R}^d$. Then $\text{VC}(\mathcal{S}) = d + 1$, $\text{vc}(\mathcal{S}) = d$.

VC density is often the right measure for the combinatorial complexity of a set system. (E.g., it is related to packing numbers and entropy).

Some basic properties:

- $\text{vc}(\mathcal{S}) \leq \text{VC}(\mathcal{S})$, and if one is finite then so is the other;
- $\text{VC}(\mathcal{S}) = 0 \iff |\mathcal{S}| = 1$;
- \mathcal{S} is finite $\iff \text{vc}(\mathcal{S}) = 0 \iff \text{vc}(\mathcal{S}) < 1$;
- $\mathcal{S} = \mathcal{S}_1 \cup \mathcal{S}_2 \Rightarrow \text{vc}(\mathcal{S}) = \max\{\text{vc}(\mathcal{S}_1), \text{vc}(\mathcal{S}_2)\}$.

Let X be a set (possibly finite). Given $A_1, \dots, A_n \subseteq X$, denote by $S(A_1, \dots, A_n)$ the set of atoms of the Boolean subalgebra of 2^X generated by A_1, \dots, A_n : those subsets of X of the form

$$\bigcap_{i \in I} A_i \cap \bigcap_{i \notin I} X \setminus A_i \quad \text{where } I \subseteq \{1, \dots, n\}$$

which are *non-empty* (= “the non-empty sets in the Venn diagram of A_1, \dots, A_n ”).

Let X be a set (possibly finite). Given $A_1, \dots, A_n \subseteq X$, denote by $S(A_1, \dots, A_n)$ the set of atoms of the Boolean subalgebra of 2^X generated by A_1, \dots, A_n : those subsets of X of the form

$$\bigcap_{i \in I} A_i \cap \bigcap_{i \notin I} X \setminus A_i \quad \text{where } I \subseteq \{1, \dots, n\}$$

which are *non-empty* (= “the non-empty sets in the Venn diagram of A_1, \dots, A_n ”).

Suppose now that \mathcal{S} is a set system on X . We define

$$n \mapsto \pi_{\mathcal{S}}^*(n) := \max \{ |S(A_1, \dots, A_n)| : A_1, \dots, A_n \in \mathcal{S} \} : \mathbb{N} \rightarrow \mathbb{N}.$$

Let X be a set (possibly finite). Given $A_1, \dots, A_n \subseteq X$, denote by $S(A_1, \dots, A_n)$ the set of atoms of the Boolean subalgebra of 2^X generated by A_1, \dots, A_n : those subsets of X of the form

$$\bigcap_{i \in I} A_i \cap \bigcap_{i \notin I} X \setminus A_i \quad \text{where } I \subseteq \{1, \dots, n\}$$

which are *non-empty* (= “the non-empty sets in the Venn diagram of A_1, \dots, A_n ”).

Suppose now that \mathcal{S} is a set system on X . We define

$$n \mapsto \pi_{\mathcal{S}}^*(n) := \max \{ |S(A_1, \dots, A_n)| : A_1, \dots, A_n \in \mathcal{S} \} : \mathbb{N} \rightarrow \mathbb{N}.$$

We say that \mathcal{S} is **independent** (in X) if $\pi_{\mathcal{S}}^*(n) = 2^n$ for every n , and **dependent** (in X) otherwise.

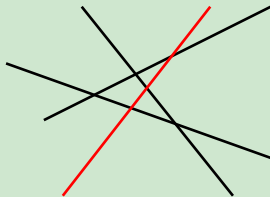
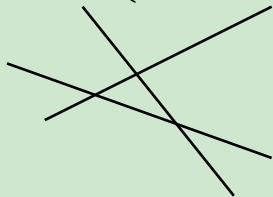
Example ($X = \mathbb{R}^2$, $\mathcal{S} =$ half planes in \mathbb{R}^2)

Example ($X = \mathbb{R}^2$, $\mathcal{S} =$ half planes in \mathbb{R}^2)

$$\pi_{\mathcal{S}}^*(n) = \begin{cases} \text{maximum number of regions into which } n \text{ half} \\ \text{planes partition the plane.} \end{cases}$$

Example ($X = \mathbb{R}^2$, $\mathcal{S} =$ half planes in \mathbb{R}^2)

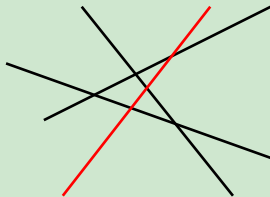
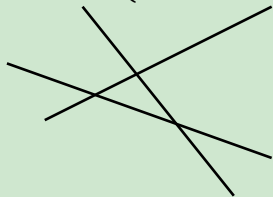
$\pi_{\mathcal{S}}^*(n) = \left\{ \begin{array}{l} \text{maximum number of regions into which } n \text{ half} \\ \text{planes partition the plane.} \end{array} \right.$



Adding one half plane to $n - 1$ given half planes divides at most n of the existing regions into 2 pieces. So $\pi_{\mathcal{S}}^*(n) = O(n^2)$.

Example ($X = \mathbb{R}^2$, $\mathcal{S} =$ half planes in \mathbb{R}^2)

$\pi_{\mathcal{S}}^*(n) = \left\{ \begin{array}{l} \text{maximum number of regions into which } n \text{ half} \\ \text{planes partition the plane.} \end{array} \right.$



Adding one half plane to $n - 1$ given half planes divides at most n of the existing regions into 2 pieces. So $\pi_{\mathcal{S}}^*(n) = O(n^2)$.

The function $\pi_{\mathcal{S}}^*$ is called the **dual shatter function of \mathcal{S}** .

Let X, Y be infinite sets, $\Phi \subseteq X \times Y$ a binary relation.

Let X, Y be infinite sets, $\Phi \subseteq X \times Y$ a binary relation. Put

$$\mathcal{S}_\Phi := \{\Phi_y : y \in Y\} \subseteq 2^X \quad \text{where } \Phi_y := \{x \in X : (x, y) \in \Phi\},$$

Let X, Y be infinite sets, $\Phi \subseteq X \times Y$ a binary relation. Put

$$\mathcal{S}_\Phi := \{\Phi_y : y \in Y\} \subseteq 2^X \quad \text{where } \Phi_y := \{x \in X : (x, y) \in \Phi\},$$

and

$$\begin{aligned} \pi_\Phi &:= \pi_{\mathcal{S}_\Phi}, & \pi_\Phi^* &:= \pi_{\mathcal{S}_\Phi}^*, \\ \text{VC}(\Phi) &:= \text{VC}(\mathcal{S}_\Phi), & \text{vc}(\Phi) &:= \text{vc}(\mathcal{S}_\Phi). \end{aligned}$$

Let X, Y be infinite sets, $\Phi \subseteq X \times Y$ a binary relation. Put

$$\mathcal{S}_\Phi := \{\Phi_y : y \in Y\} \subseteq 2^X \quad \text{where } \Phi_y := \{x \in X : (x, y) \in \Phi\},$$

and

$$\begin{aligned} \pi_\Phi &:= \pi_{\mathcal{S}_\Phi}, & \pi_\Phi^* &:= \pi_{\mathcal{S}_\Phi}^*, \\ \text{VC}(\Phi) &:= \text{VC}(\mathcal{S}_\Phi), & \text{vc}(\Phi) &:= \text{vc}(\mathcal{S}_\Phi). \end{aligned}$$

We also write

$$\Phi^* \subseteq Y \times X := \{(y, x) \in Y \times X : (x, y) \in \Phi\}.$$

Let X, Y be infinite sets, $\Phi \subseteq X \times Y$ a binary relation. Put

$$\mathcal{S}_\Phi := \{\Phi_y : y \in Y\} \subseteq 2^X \quad \text{where } \Phi_y := \{x \in X : (x, y) \in \Phi\},$$

and

$$\begin{aligned} \pi_\Phi &:= \pi_{\mathcal{S}_\Phi}, & \pi_\Phi^* &:= \pi_{\mathcal{S}_\Phi}^*, \\ \text{VC}(\Phi) &:= \text{VC}(\mathcal{S}_\Phi), & \text{vc}(\Phi) &:= \text{vc}(\mathcal{S}_\Phi). \end{aligned}$$

We also write

$$\Phi^* \subseteq Y \times X := \{(y, x) \in Y \times X : (x, y) \in \Phi\}.$$

In this way we obtain two set systems: (X, \mathcal{S}_Φ) and $(Y, \mathcal{S}_{\Phi^*})$

Let X, Y be infinite sets, $\Phi \subseteq X \times Y$ a binary relation. Put

$$\mathcal{S}_\Phi := \{\Phi_y : y \in Y\} \subseteq 2^X \quad \text{where } \Phi_y := \{x \in X : (x, y) \in \Phi\},$$

and

$$\begin{aligned} \pi_\Phi &:= \pi_{\mathcal{S}_\Phi}, & \pi_\Phi^* &:= \pi_{\mathcal{S}_\Phi^*}, \\ \text{VC}(\Phi) &:= \text{VC}(\mathcal{S}_\Phi), & \text{vc}(\Phi) &:= \text{vc}(\mathcal{S}_\Phi). \end{aligned}$$

We also write

$$\Phi^* \subseteq Y \times X := \{(y, x) \in Y \times X : (x, y) \in \Phi\}.$$

In this way we obtain two set systems: (X, \mathcal{S}_Φ) and $(Y, \mathcal{S}_{\Phi^*})$

Given a finite set $A \subseteq X$ we have a bijection

$$B \mapsto \bigcap_{x \in B} \Phi_x^* \cap \bigcap_{x \in A \setminus B} Y \setminus \Phi_x^* : \quad \mathcal{S}_\Phi \cap A \rightarrow \mathcal{S}(\Phi_x^* : x \in A).$$

Hence $\pi_{\Phi} = \pi_{\Phi^*}^*$ and $\pi_{\Phi^*} = \pi_{\Phi}^*$, and thus

$$\begin{aligned} \mathcal{S}_{\Phi} \text{ is a VC class} &\iff \mathcal{S}_{\Phi^*} \text{ is dependent,} \\ \mathcal{S}_{\Phi^*} \text{ is a VC class} &\iff \mathcal{S}_{\Phi} \text{ is dependent.} \end{aligned}$$

Hence $\pi_{\Phi} = \pi_{\Phi^*}^*$ and $\pi_{\Phi^*} = \pi_{\Phi}^*$, and thus

$$\begin{aligned} \mathcal{S}_{\Phi} \text{ is a VC class} &\iff \mathcal{S}_{\Phi^*} \text{ is dependent,} \\ \mathcal{S}_{\Phi^*} \text{ is a VC class} &\iff \mathcal{S}_{\Phi} \text{ is dependent.} \end{aligned}$$

Moreover (first noticed by Assouad):

$$\mathcal{S}_{\Phi} \text{ is a VC class} \iff \mathcal{S}_{\Phi^*} \text{ is a VC class.}$$

The model-theoretic context

We fix:

\mathcal{L} : a first-order language,

$x = (x_1, \dots, x_m)$: object variables,

$y = (y_1, \dots, y_n)$: parameter variables,

$\varphi(x; y)$: a partitioned \mathcal{L} -formula,

M : an infinite \mathcal{L} -structure, and

T : a complete \mathcal{L} -theory without finite models.

We fix:

\mathcal{L} : a first-order language,

$x = (x_1, \dots, x_m)$: object variables,

$y = (y_1, \dots, y_n)$: parameter variables,

$\varphi(x; y)$: a partitioned \mathcal{L} -formula,

M : an infinite \mathcal{L} -structure, and

T : a complete \mathcal{L} -theory without finite models.

The set system (on M^m) associated with φ in M :

$$\mathcal{S}_\varphi^M := \{\varphi^M(M^m; b) : b \in M^n\}$$

We fix:

\mathcal{L} : a first-order language,

$x = (x_1, \dots, x_m)$: object variables,

$y = (y_1, \dots, y_n)$: parameter variables,

$\varphi(x; y)$: a partitioned \mathcal{L} -formula,

M : an infinite \mathcal{L} -structure, and

T : a complete \mathcal{L} -theory without finite models.

The set system (on M^m) associated with φ in M :

$$\mathcal{S}_\varphi^M := \{\varphi^M(M^m; b) : b \in M^n\}$$

If $M \equiv N$, then $\pi_{\mathcal{S}_\varphi^M} = \pi_{\mathcal{S}_\varphi^N}$. So, picking $M \models T$ arbitrary, set

$$\pi_\varphi := \pi_{\mathcal{S}_\varphi^M}, \quad \text{VC}(\varphi) := \text{VC}(\mathcal{S}_\varphi^M), \quad \text{vc}(\varphi) := \text{vc}(\mathcal{S}_\varphi^M).$$

The *dual* of $\varphi(x; y)$ is $\varphi^*(y; x) := \varphi(x; y)$. Put

$$\text{VC}^*(\varphi) := \text{VC}(\varphi^*), \quad \text{vc}^*(\varphi) := \text{vc}(\varphi^*).$$

The *dual* of $\varphi(x; y)$ is $\varphi^*(y; x) := \varphi(x; y)$. Put

$$\text{VC}^*(\varphi) := \text{VC}(\varphi^*), \quad \text{vc}^*(\varphi) := \text{vc}(\varphi^*).$$

We have $\pi_\varphi^* = \pi_{\varphi^*}$, hence $\text{VC}^*(\varphi)$ and $\text{vc}^*(\varphi)$ can be computed using the dual shatter function of φ .

The *dual* of $\varphi(x; y)$ is $\varphi^*(y; x) := \varphi(x; y)$. Put

$$\text{VC}^*(\varphi) := \text{VC}(\varphi^*), \quad \text{vc}^*(\varphi) := \text{vc}(\varphi^*).$$

We have $\pi_\varphi^* = \pi_{\varphi^*}$, hence $\text{VC}^*(\varphi)$ and $\text{vc}^*(\varphi)$ can be computed using the dual shatter function of φ .

Recall:

If $\text{VC}(\varphi) < \infty$ then we say that φ is **dependent** in T . The theory T does **not have the independence property** (is **NIP**, or **dependent**) if every partitioned \mathcal{L} -formula is dependent in T .

The *dual* of $\varphi(x; y)$ is $\varphi^*(y; x) := \varphi(x; y)$. Put

$$\text{VC}^*(\varphi) := \text{VC}(\varphi^*), \quad \text{vc}^*(\varphi) := \text{vc}(\varphi^*).$$

We have $\pi_\varphi^* = \pi_{\varphi^*}$, hence $\text{VC}^*(\varphi)$ and $\text{vc}^*(\varphi)$ can be computed using the dual shatter function of φ .

Recall:

If $\text{VC}(\varphi) < \infty$ then we say that φ is **dependent** in T . The theory T does **not have the independence property** (is **NIP**, or **dependent**) if every partitioned \mathcal{L} -formula is dependent in T .

An important theorem of Shelah (given other proofs by Laskowski and others) says that for T to be NIP it is enough for for every \mathcal{L} -formula $\varphi(x; y)$ with $|x| = 1$ to be dependent.

Some questions about vc in model theory

Some questions about vc in model theory

- 1 Possible values of $vc(\varphi)$.

Some questions about vc in model theory

- 1 Possible values of $\text{vc}(\varphi)$. There exists a formula $\varphi(x; y)$ in $\mathcal{L}_{\text{rings}}$ with $|y| = 4$ such that

$$\text{vc}^{\text{ACF}_0}(\varphi) = \frac{4}{3}; \quad \text{vc}^{\text{ACF}_p}(\varphi) = \frac{3}{2} \text{ for } p > 0.$$

Some questions about vc in model theory

- 1 Possible values of $\text{vc}(\varphi)$. There exists a formula $\varphi(x; y)$ in $\mathcal{L}_{\text{rings}}$ with $|y| = 4$ such that

$$\text{vc}^{\text{ACF}_0}(\varphi) = \frac{4}{3}; \quad \text{vc}^{\text{ACF}_p}(\varphi) = \frac{3}{2} \text{ for } p > 0.$$

We do not know an example of a formula φ in a NIP theory with $\text{vc}(\varphi) \notin \mathbb{Q}$.

Some questions about vc in model theory

- 1 Possible values of $\text{vc}(\varphi)$. There exists a formula $\varphi(x; y)$ in $\mathcal{L}_{\text{rings}}$ with $|y| = 4$ such that

$$\text{vc}^{\text{ACF}_0}(\varphi) = \frac{4}{3}; \quad \text{vc}^{\text{ACF}_p}(\varphi) = \frac{3}{2} \text{ for } p > 0.$$

We do not know an example of a formula φ in a NIP theory with $\text{vc}(\varphi) \notin \mathbb{Q}$.

- 2 Growth of π_φ .

Some questions about vc in model theory

- ① Possible values of $\text{vc}(\varphi)$. There exists a formula $\varphi(x; y)$ in $\mathcal{L}_{\text{rings}}$ with $|y| = 4$ such that

$$\text{vc}^{\text{ACF}_0}(\varphi) = \frac{4}{3}; \quad \text{vc}^{\text{ACF}_p}(\varphi) = \frac{3}{2} \text{ for } p > 0.$$

We do not know an example of a formula φ in a NIP theory with $\text{vc}(\varphi) \notin \mathbb{Q}$.

- ② Growth of π_φ . There is an example of an ω -stable T and an \mathcal{L} -formula $\varphi(x; y)$ with $|y| = 2$ and

$$\pi_\varphi(n) = \frac{1}{2}n \log n (1 + o(1)).$$

- ③ Uniform bounds on $\text{vc}(\varphi)$.

Some reasons why it should be interesting to obtain bounds on $vc(\varphi)$ in terms of $|y| = \text{number of free parameters}$:

Some reasons why it should be interesting to obtain bounds on $vc(\varphi)$ in terms of $|y| =$ number of free parameters:

- 1 uniform bounds on VC density often “explain” why certain bounds on the complexity of geometric arrangements, used in computational geometry, are polynomial in the number of objects involved (*example follows later*);

Some reasons why it should be interesting to obtain bounds on $vc(\varphi)$ in terms of $|y| =$ number of free parameters:

- 1 uniform bounds on VC density often “explain” why certain bounds on the complexity of geometric arrangements, used in computational geometry, are polynomial in the number of objects involved (*example follows later*);
- 2 connections to strengthenings of the NIP concept: if $vc(\varphi) < 2$ for each $\varphi(x; y)$ with $|y| = 1$, then T is *dp-minimal*.

Theorem

*Suppose T expands the theory of linearly ordered sets, and assume that T is **weakly o-minimal**, i.e., in every $M \models T$, every definable subset of M is a finite union of convex subsets of M .*

Theorem

*Suppose T expands the theory of linearly ordered sets, and assume that T is **weakly o-minimal**, i.e., in every $M \models T$, every definable subset of M is a finite union of convex subsets of M . Then for each $\varphi(x; y)$ we have $\pi_\varphi(t) = O(t^{|y|})$, hence $\text{vc}(\varphi) \leq |y|$.*

Theorem

*Suppose T expands the theory of linearly ordered sets, and assume that T is **weakly o-minimal**, i.e., in every $M \models T$, every definable subset of M is a finite union of convex subsets of M . Then for each $\varphi(x; y)$ we have $\pi_\varphi(t) = O(t^{|y|})$, hence $\text{vc}(\varphi) \leq |y|$.*

(Generalizes earlier results due to Karpinski-Macintyre and Wilkie.)

Theorem

*Suppose T expands the theory of linearly ordered sets, and assume that T is **weakly o-minimal**, i.e., in every $M \models T$, every definable subset of M is a finite union of convex subsets of M . Then for each $\varphi(x; y)$ we have $\pi_\varphi(t) = O(t^{|y|})$, hence $\text{vc}(\varphi) \leq |y|$.*

(Generalizes earlier results due to Karpinski-Macintyre and Wilkie.)

We sketch the proof.

Theorem

*Suppose T expands the theory of linearly ordered sets, and assume that T is **weakly o-minimal**, i.e., in every $M \models T$, every definable subset of M is a finite union of convex subsets of M . Then for each $\varphi(x; y)$ we have $\pi_\varphi(t) = O(t^{|y|})$, hence $\text{vc}(\varphi) \leq |y|$.*

(Generalizes earlier results due to Karpinski-Macintyre and Wilkie.)

We sketch the proof.

It is more convenient to work with π^* , and thus we need to show

$$\pi_\varphi^*(t) = O(t^{|x|}) \quad \text{for each } \varphi(x; y).$$

It is also convenient to be able to deal with finitely many formulas at once:

It is also convenient to be able to deal with finitely many formulas at once:

$\Delta(x; y)$: a finite non-empty set of partitioned \mathcal{L} -formulas;

$S^\Delta(B)$: the set of complete $\Delta(x; B)$ -types in M ($B \subseteq M^{|y|}$).

It is also convenient to be able to deal with finitely many formulas at once:

$\Delta(x; y)$: a finite non-empty set of partitioned \mathcal{L} -formulas;

$S^\Delta(B)$: the set of complete $\Delta(x; B)$ -types in M ($B \subseteq M^{|y|}$).

If T is NIP then we set

$$\pi_\Delta^*(t) := \max \left\{ |S^\Delta(B)| : B \in \binom{M^{|y|}}{t} \right\},$$
$$\text{vc}^*(\Delta) := \inf \{ r \in \mathbb{R}^{>0} : \pi_\Delta^*(t) = O(t^r) \}.$$

Definition (adapted from Guingona)

Δ has **uniform definability of types over finite sets (UDTFS)** in M with m parameters if there are families of \mathcal{L} -formulas

$$\mathcal{F}_i = (\varphi_i(y; y_1, \dots, y_m))_{\varphi \in \Delta} \quad (i \in I = \text{a finite set})$$

such that for every finite $B \subseteq M^{|y|}$ and $q \in S^\Delta(B)$ there are $b_1, \dots, b_m \in B$ and $i \in I$ such that $\mathcal{F}_i(y; b_1, \dots, b_m)$ defines q .

Definition (adapted from Guingona)

Δ has **uniform definability of types over finite sets (UDTFS)** in M with m parameters if there are families of \mathcal{L} -formulas

$$\mathcal{F}_i = (\varphi_i(y; y_1, \dots, y_m))_{\varphi \in \Delta} \quad (i \in I = \text{a finite set})$$

such that for every finite $B \subseteq M^{|y|}$ and $q \in S^\Delta(B)$ there are $b_1, \dots, b_m \in B$ and $i \in I$ such that $\mathcal{F}_i(y; b_1, \dots, b_m)$ defines q .

- If we don't care about the number of extra parameters m , then we can always achieve $|I| = 1$ and $|\Delta| = 1$.

Definition (adapted from Guingona)

Δ has **uniform definability of types over finite sets (UDTFS)** in M with m parameters if there are families of \mathcal{L} -formulas

$$\mathcal{F}_i = (\varphi_i(y; y_1, \dots, y_m))_{\varphi \in \Delta} \quad (i \in I = \text{a finite set})$$

such that for every finite $B \subseteq M^{|y|}$ and $q \in S^\Delta(B)$ there are $b_1, \dots, b_m \in B$ and $i \in I$ such that $\mathcal{F}_i(y; b_1, \dots, b_m)$ defines q .

- If we don't care about the number of extra parameters m , then we can always achieve $|I| = 1$ and $|\Delta| = 1$.
- On the other hand: if Δ has a uniform definition $\mathcal{F} = (\mathcal{F}_i)_{i \in I}$ for Δ -types with m parameters, then

$$|S^\Delta(B)| \leq |I| \cdot |B|^m \quad \text{for every finite } B.$$

Theorem

*Suppose that M has the VC m **property**, i.e., any $\Delta(x; y)$ with $|x| = 1$ has UDTFS in M with m parameters.*

Theorem

*Suppose that M has the VC m **property**, i.e., any $\Delta(x; y)$ with $|x| = 1$ has UDTFS in M with m parameters.*

Then every $\Delta(x; y)$ has UDTFS in M with $m|x|$ parameters.

Weakly o-minimal theories have the VC 1 property (sketch).

Weakly o-minimal theories have the VC 1 property (sketch).

Let $M \models T$ and $\Delta(x; y)$ be a finite non-empty set of \mathcal{L} -formulas with $|x| = 1$. We let φ range over Δ and b over $M^{|y|}$.

Weakly o-minimal theories have the VC 1 property (sketch).

Let $M \models T$ and $\Delta(x; y)$ be a finite non-empty set of \mathcal{L} -formulas with $|x| = 1$. We let φ range over Δ and b over $M^{|y|}$.

If for each φ and b , the set $\varphi(M; b)$ is an initial segment of M , then clearly Δ has UDTFS with a single parameter.

Weakly o-minimal theories have the VC 1 property (sketch).

Let $M \models T$ and $\Delta(x; y)$ be a finite non-empty set of \mathcal{L} -formulas with $|x| = 1$. We let φ range over Δ and b over $M^{|y|}$.

If for each φ and b , the set $\varphi(M; b)$ is an initial segment of M , then clearly Δ has UDTFS with a single parameter.

In general, there is some N such that for each φ and b , $\varphi(M; b)$ has $\leq N$ convex components, and hence is a Boolean combination of $\leq 2N$ initial segments of M (uniformly in b).

Weakly o-minimal theories have the VC 1 property (sketch).

Let $M \models T$ and $\Delta(x; y)$ be a finite non-empty set of \mathcal{L} -formulas with $|x| = 1$. We let φ range over Δ and b over $M^{|y|}$.

If for each φ and b , the set $\varphi(M; b)$ is an initial segment of M , then clearly Δ has UDTFS with a single parameter.

In general, there is some N such that for each φ and b , $\varphi(M; b)$ has $\leq N$ convex components, and hence is a Boolean combination of $\leq 2N$ initial segments of M (uniformly in b).

Forming Boolean combinations preserves UDTFS. □

Weakly o-minimal theories have the VC 1 property (sketch).

Let $M \models T$ and $\Delta(x; y)$ be a finite non-empty set of \mathcal{L} -formulas with $|x| = 1$. We let φ range over Δ and b over $M^{|y|}$.

If for each φ and b , the set $\varphi(M; b)$ is an initial segment of M , then clearly Δ has UDTFS with a single parameter.

In general, there is some N such that for each φ and b , $\varphi(M; b)$ has $\leq N$ convex components, and hence is a Boolean combination of $\leq 2N$ initial segments of M (uniformly in b).

Forming Boolean combinations preserves UDTFS. □

The same proof applies to quasi-o-minimal theories (e.g., Presburger Arithmetic).

Uniform bounds on VC density

Interesting classes of NIP theories are provided by certain valued fields.

Interesting classes of NIP theories are provided by certain valued fields. By a non-trivial elaboration of our methods:

Theorem

Suppose $M = \mathbb{Q}_p$ is the field of p -adic numbers, construed as a structure in the language of rings.

Interesting classes of NIP theories are provided by certain valued fields. By a non-trivial elaboration of our methods:

Theorem

Suppose $M = \mathbb{Q}_p$ is the field of p -adic numbers, construed as a structure in the language of rings. Then M has the VC 2 property; in fact, $\text{vc}(\varphi) \leq 2|y| - 1$.

Interesting classes of NIP theories are provided by certain valued fields. By a non-trivial elaboration of our methods:

Theorem

Suppose $M = \mathbb{Q}_p$ is the field of p -adic numbers, construed as a structure in the language of rings. Then M has the VC 2 property; in fact, $\text{vc}(\varphi) \leq 2|y| - 1$.

This (probably non-optimal) result also holds, e.g., for the subanalytic expansions of \mathbb{Q}_p considered by Denef & v. d. Dries.

Interesting classes of NIP theories are provided by certain valued fields. By a non-trivial elaboration of our methods:

Theorem

Suppose $M = \mathbb{Q}_p$ is the field of p -adic numbers, construed as a structure in the language of rings. Then M has the VC 2 property; in fact, $\text{vc}(\varphi) \leq 2|y| - 1$.

This (probably non-optimal) result also holds, e.g., for the subanalytic expansions of \mathbb{Q}_p considered by Denef & v. d. Dries.

We do not know whether the completions of ACVF have the VC d property.

Interesting classes of NIP theories are provided by certain valued fields. By a non-trivial elaboration of our methods:

Theorem

Suppose $M = \mathbb{Q}_p$ is the field of p -adic numbers, construed as a structure in the language of rings. Then M has the VC 2 property; in fact, $\text{vc}(\varphi) \leq 2|y| - 1$.

This (probably non-optimal) result also holds, e.g., for the subanalytic expansions of \mathbb{Q}_p considered by Denef & v. d. Dries.

We do not know whether the completions of ACVF have the VC d property.

We also have results stating that in certain stable theories T we have linear bounds on VC density, not obtained via the VC m property.

Theorem

Let A be an infinite abelian group. T.f.a.e.:

- 1 $\text{vc}(\varphi)$ for $\varphi(x; y)$ with $|y| = 1$ is bounded;
- 2 there is some d such that $\text{vc}(\varphi) \leq d|y|$ for each $\varphi(x; y)$;
- 3 there are only finitely many p such that $A[p]$ or A/pA is infinite, and for all p there are only finitely many n such that

$$U(p, n; A) = |(p^n A)[p]/(p^{n+1} A)[p]| \geq \aleph_0.$$

Theorem

Let A be an infinite abelian group. T.f.a.e.:

- 1 $\text{vc}(\varphi)$ for $\varphi(x; y)$ with $|y| = 1$ is bounded;
- 2 there is some d such that $\text{vc}(\varphi) \leq d|y|$ for each $\varphi(x; y)$;
- 3 there are only finitely many p such that $A[p]$ or A/pA is infinite, and for all p there are only finitely many n such that

$$U(p, n; A) = |(p^n A)[p]/(p^{n+1} A)[p]| \geq \aleph_0.$$

As an upshot of the proof of the theorem we are able to determine the theories of all dp-minimal abelian groups.

Theorem

Let A be an infinite abelian group. T.f.a.e.:

- 1 $\text{vc}(\varphi)$ for $\varphi(x; y)$ with $|y| = 1$ is bounded;
- 2 there is some d such that $\text{vc}(\varphi) \leq d|y|$ for each $\varphi(x; y)$;
- 3 there are only finitely many p such that $A[p]$ or A/pA is infinite, and for all p there are only finitely many n such that

$$U(p, n; A) = |(p^n A)[p]/(p^{n+1} A)[p]| \geq \aleph_0.$$

As an upshot of the proof of the theorem we are able to determine the theories of all dp-minimal abelian groups.

If A has finite exponent then it has the VC d property (explicit d). The proof involves some combinatorics with distributive lattices.

A general theorem is:

Theorem

Suppose T does not have the finite cover property and finite U -rank $U(T)$. Then $\text{vc}(\varphi) \leq |y| U(T)$ for every $\varphi(x; y)$.

A general theorem is:

Theorem

Suppose T does not have the finite cover property and finite U -rank $U(T)$. Then $\text{vc}(\varphi) \leq |y| U(T)$ for every $\varphi(x; y)$.

Cases where the theorem applies includes all expansions T of the theory of groups with $\text{MR}(T) < \omega$.

A general theorem is:

Theorem

Suppose T does not have the finite cover property and finite U -rank $U(T)$. Then $\text{vc}(\varphi) \leq |y| U(T)$ for every $\varphi(x; y)$.

Cases where the theorem applies includes all expansions T of the theory of groups with $\text{MR}(T) < \omega$. The theorem in action:

A general theorem is:

Theorem

Suppose T does not have the finite cover property and finite U-rank $U(T)$. Then $vc(\varphi) \leq |y| U(T)$ for every $\varphi(x; y)$.

Cases where the theorem applies includes all expansions T of the theory of groups with $MR(T) < \omega$. The theorem in action:

Example (\mathcal{L} = language of rings, $K \models \text{ACF}$)

Choose $\varphi(x; y)$ so that $\mathcal{S}_\varphi^K =$ all zero sets (in K^m) of polynomials in m indeterminates over K of degree $\leq d$.

A general theorem is:

Theorem

Suppose T does not have the finite cover property and finite U -rank $U(T)$. Then $\text{vc}(\varphi) \leq |y| U(T)$ for every $\varphi(x; y)$.

Cases where the theorem applies includes all expansions T of the theory of groups with $\text{MR}(T) < \omega$. The theorem in action:

Example (\mathcal{L} = language of rings, $K \models \text{ACF}$)

Choose $\varphi(x; y)$ so that \mathcal{S}_φ^K = all zero sets (in K^m) of polynomials in m indeterminates over K of degree $\leq d$. Then

$$\pi_\varphi^*(t) = \left\{ \begin{array}{l} \text{maximum number of non-empty} \\ \text{Boolean combinations of } t \text{ hypersur-} \\ \text{faces in } K^m \text{ of degree } \leq d. \end{array} \right\} = \pi_{\varphi^*}(t) = O(t^m)$$

Question

Let $f: A \rightarrow \mathbb{R}^n$, $A \subseteq \mathbb{R}^m$, be L -Lipschitz (where $L \in \mathbb{R}^{\geq 0}$), i.e.,

$$\|f(x) - f(y)\| \leq L \cdot \|x - y\| \quad \text{for all } x, y \in A.$$

Can one extend f to an L -Lipschitz map $\mathbb{R}^m \rightarrow \mathbb{R}^n$?

Question

Let $f: A \rightarrow \mathbb{R}^n$, $A \subseteq \mathbb{R}^m$, be L -Lipschitz (where $L \in \mathbb{R}^{\geq 0}$), i.e.,

$$\|f(x) - f(y)\| \leq L \cdot \|x - y\| \quad \text{for all } x, y \in A.$$

Can one extend f to an L -Lipschitz map $\mathbb{R}^m \rightarrow \mathbb{R}^n$?

Kirszbraun (1934): yes for all n

There always exists an L -Lipschitz extension $\mathbb{R}^m \rightarrow \mathbb{R}^n$ of f .

Question

Let $f: A \rightarrow \mathbb{R}^n$, $A \subseteq \mathbb{R}^m$, be L -Lipschitz (where $L \in \mathbb{R}^{\geq 0}$), i.e.,

$$\|f(x) - f(y)\| \leq L \cdot \|x - y\| \quad \text{for all } x, y \in A.$$

Can one extend f to an L -Lipschitz map $\mathbb{R}^m \rightarrow \mathbb{R}^n$?

Kirszbraun (1934): yes for all n

There always exists an L -Lipschitz extension $\mathbb{R}^m \rightarrow \mathbb{R}^n$ of f .

The usual proofs of this theorem all use some sort of transfinite induction. (A classical explicit construction by MacShane & Whitney only yields an $L\sqrt{n}$ -Lipschitz extension.)

Theorem A (A.-Fischer, Proc. LMS 2011)

Let $\mathcal{R} = (R, 0, 1, +, \times, <, \dots)$ be a **definably complete** expansion of an ordered field: every non-empty definable subset of R which is bounded from above has a supremum.

Theorem A (A.-Fischer, Proc. LMS 2011)

Let $\mathcal{R} = (R, 0, 1, +, \times, <, \dots)$ be a **definably complete** expansion of an ordered field: every non-empty definable subset of R which is bounded from above has a supremum. Then every definable L -Lipschitz map $A \rightarrow R^n$ ($A \subseteq R^m$, $L \in R^{\geq 0}$) has a definable L -Lipschitz extension $R^m \rightarrow R^n$.

Theorem A (A.-Fischer, Proc. LMS 2011)

Let $\mathcal{R} = (R, 0, 1, +, \times, <, \dots)$ be a **definably complete** expansion of an ordered field: every non-empty definable subset of R which is bounded from above has a supremum. Then every definable L -Lipschitz map $A \rightarrow R^n$ ($A \subseteq R^m$, $L \in R^{\geq 0}$) has a definable L -Lipschitz extension $R^m \rightarrow R^n$.

The proof of this theorem used convex analysis and is based on a relationship between Lipschitz maps and monotone set-valued maps (Minty; more recently, Bauschke & Wang).

Theorem A (A.-Fischer, Proc. LMS 2011)

Let $\mathcal{R} = (R, 0, 1, +, \times, <, \dots)$ be a **definably complete** expansion of an ordered field: every non-empty definable subset of R which is bounded from above has a supremum. Then every definable L -Lipschitz map $A \rightarrow R^n$ ($A \subseteq R^m$, $L \in R^{\geq 0}$) has a definable L -Lipschitz extension $R^m \rightarrow R^n$.

The proof of this theorem used convex analysis and is based on a relationship between Lipschitz maps and monotone set-valued maps (Minty; more recently, Bauschke & Wang).

Another crucial ingredient (in the case where $R \neq \mathbb{R}$) is a definable version of a classical theorem of Helly:

Theorem B (A.-Fischer, Proc. LMS 2011)

Let R be a definably complete expansion of an ordered field.
Let \mathcal{C} be a definable family of closed bounded *convex* subsets of R^n .

Theorem B (A.-Fischer, Proc. LMS 2011)

Let R be a definably complete expansion of an ordered field.
Let \mathcal{C} be a definable family of closed bounded *convex* subsets of R^n . Suppose \mathcal{C} is $(n + 1)$ -**consistent**:

$$\bigcap C' \neq \emptyset \quad \text{for all } C' \subseteq \mathcal{C} \text{ with } |C'| \leq n + 1.$$

Theorem B (A.-Fischer, Proc. LMS 2011)

Let R be a definably complete expansion of an ordered field.
Let \mathcal{C} be a definable family of closed bounded *convex* subsets of R^n . Suppose \mathcal{C} is $(n + 1)$ -**consistent**:

$$\bigcap \mathcal{C}' \neq \emptyset \quad \text{for all } \mathcal{C}' \subseteq \mathcal{C} \text{ with } |\mathcal{C}'| \leq n + 1.$$

Then $\bigcap \mathcal{C} \neq \emptyset$.

Theorem B (A.-Fischer, Proc. LMS 2011)

Let R be a definably complete expansion of an ordered field.
Let \mathcal{C} be a definable family of closed bounded *convex* subsets of R^n . Suppose \mathcal{C} is $(n + 1)$ -**consistent**:

$$\bigcap \mathcal{C}' \neq \emptyset \quad \text{for all } \mathcal{C}' \subseteq \mathcal{C} \text{ with } |\mathcal{C}'| \leq n + 1.$$

Then $\bigcap \mathcal{C} \neq \emptyset$.

Our proof of this theorem uses an optimization argument.

Theorem B (A.-Fischer, Proc. LMS 2011)

Let R be a definably complete expansion of an ordered field. Let \mathcal{C} be a definable family of closed bounded *convex* subsets of R^n . Suppose \mathcal{C} is $(n + 1)$ -**consistent**:

$$\bigcap \mathcal{C}' \neq \emptyset \quad \text{for all } \mathcal{C}' \subseteq \mathcal{C} \text{ with } |\mathcal{C}'| \leq n + 1.$$

Then $\bigcap \mathcal{C} \neq \emptyset$.

Our proof of this theorem uses an optimization argument.

S. Starchenko pointed out that in the case of an o-minimal R , our theorem follows from an analysis of forking in o-minimal theories due to A. Dolich.

A subset T of X is called a **transversal** of a set system \mathcal{S} on X if every member of \mathcal{S} contains an element of T .

A subset T of X is called a **transversal** of a set system \mathcal{S} on X if every member of \mathcal{S} contains an element of T .

Theorem (Dolich '04, made explicit by Peterzil & Pillay '07)

Let \mathcal{R} be an o-minimal expansion of a real closed field, and let $\mathcal{C} = \{C_a\}_{a \in A}$ be a definable family of closed and bounded subsets of \mathcal{R}^n parameterized by a subset A of \mathcal{R}^m .

A subset T of X is called a **transversal** of a set system \mathcal{S} on X if every member of \mathcal{S} contains an element of T .

Theorem (Dolich '04, made explicit by Peterzil & Pillay '07)

Let \mathcal{R} be an o-minimal expansion of a real closed field, and let $\mathcal{C} = \{C_a\}_{a \in A}$ be a definable family of closed and bounded subsets of \mathcal{R}^n parameterized by a subset A of \mathcal{R}^m . If \mathcal{C} is $N(m, n)$ -consistent, where

$$N(m, n) = (1 + 2^m) \cdot (1 + 2^{2^m}) \cdots \quad (n \text{ factors}),$$

then \mathcal{C} has a finite transversal.

Question

Can one do better than the bound $N(m, n)$?

Question

Can one do better than the bound $N(m, n)$?

Theorem (Matoušek, 2004)

Let (X, S) be a set system of finite dual VC density $\text{vc}^(S)$.*

Question

Can one do better than the bound $N(m, n)$?

Theorem (Matoušek, 2004)

Let (X, \mathcal{S}) be a set system of finite dual VC density $\text{vc}^(\mathcal{S})$.
Suppose \mathcal{S} is d -consistent, where $d > \text{vc}^*(\mathcal{S})$.*

Question

Can one do better than the bound $N(m, n)$?

Theorem (Matoušek, 2004)

Let (X, \mathcal{S}) be a set system of finite dual VC density $\text{vc}^(\mathcal{S})$. Suppose \mathcal{S} is d -consistent, where $d > \text{vc}^*(\mathcal{S})$. Assume that X comes equipped with a topology making all sets in \mathcal{S} compact.*

Question

Can one do better than the bound $N(m, n)$?

Theorem (Matoušek, 2004)

Let (X, \mathcal{S}) be a set system of finite dual VC density $\text{vc}^(\mathcal{S})$. Suppose \mathcal{S} is d -consistent, where $d > \text{vc}^*(\mathcal{S})$. Assume that X comes equipped with a topology making all sets in \mathcal{S} compact. Then \mathcal{S} has a finite transversal.*

Question

Can one do better than the bound $N(m, n)$?

Theorem (Matoušek, 2004)

Let (X, \mathcal{S}) be a set system of finite dual VC density $\text{vc}^(\mathcal{S})$. Suppose \mathcal{S} is d -consistent, where $d > \text{vc}^*(\mathcal{S})$. Assume that X comes equipped with a topology making all sets in \mathcal{S} compact. Then \mathcal{S} has a finite transversal.*

Corollary

Let $\mathcal{C} = \{C_a\}_{a \in A}$ be a family of compact subsets of \mathbb{R}^n definable in an o-minimal structure on \mathbb{R} . If \mathcal{C} is $(n + 1)$ -consistent, then \mathcal{C} has a finite transversal.

Proof of Theorem B in the o-minimal case (Starchenko)

Proof of Theorem B in the o-minimal case (Starchenko)

Suppose \mathcal{R} is o-minimal, and write $\mathcal{C} = \{C_a\}_{a \in A}$.

Proof of Theorem B in the o-minimal case (Starchenko)

Suppose R is o-minimal, and write $\mathcal{C} = \{C_a\}_{a \in A}$.

By Helly's Theorem for finite families, the (definable) family whose members are the intersections of $n + 1$ members of \mathcal{C} is finitely consistent.

Proof of Theorem B in the o-minimal case (Starchenko)

Suppose R is o-minimal, and write $\mathcal{C} = \{C_a\}_{a \in A}$.

By Helly's Theorem for finite families, the (definable) family whose members are the intersections of $n + 1$ members of \mathcal{C} is finitely consistent.

Apply Dolich's Theorem to this family to obtain a finite set $P \subseteq R^n$ with $P \cap C_{a_1} \cap \cdots \cap C_{a_{n+1}} \neq \emptyset$ for all $a_1, \dots, a_{n+1} \in A$.

Proof of Theorem B in the o-minimal case (Starchenko)

Suppose R is o-minimal, and write $\mathcal{C} = \{C_a\}_{a \in A}$.

By Helly's Theorem for finite families, the (definable) family whose members are the intersections of $n + 1$ members of \mathcal{C} is finitely consistent.

Apply Dolich's Theorem to this family to obtain a finite set $P \subseteq R^n$ with $P \cap C_{a_1} \cap \cdots \cap C_{a_{n+1}} \neq \emptyset$ for all $a_1, \dots, a_{n+1} \in A$.

Thus

$$\mathcal{P} = \{\text{conv}(C_a \cap P)\}_{a \in A}$$

is a family of convex subsets of R^n with only finitely many distinct members, and \mathcal{P} is $(n + 1)$ -consistent.

Proof of Theorem B in the o-minimal case (Starchenko)

Suppose R is o-minimal, and write $\mathcal{C} = \{C_a\}_{a \in A}$.

By Helly's Theorem for finite families, the (definable) family whose members are the intersections of $n + 1$ members of \mathcal{C} is finitely consistent.

Apply Dolich's Theorem to this family to obtain a finite set $P \subseteq R^n$ with $P \cap C_{a_1} \cap \cdots \cap C_{a_{n+1}} \neq \emptyset$ for all $a_1, \dots, a_{n+1} \in A$.

Thus

$$\mathcal{P} = \{\text{conv}(C_a \cap P)\}_{a \in A}$$

is a family of convex subsets of R^n with only finitely many distinct members, and \mathcal{P} is $(n + 1)$ -consistent.

Hence $\emptyset \neq \bigcap \mathcal{P} \subseteq \bigcap \mathcal{C}$ by Helly's Theorem for finite families. \square

There are many open questions in this subject. Here is one:

There are many open questions in this subject. Here is one:

Open question

Suppose T is a NIP theory.

If there is some d_1 such that $\text{vc}(\varphi) \leq d_1$ for each $\varphi(x; y)$ with $|y| = 1$, is there is some d_m such that $\text{vc}(\varphi) \leq d_m$ for each $\varphi(x; y)$ with $|y| = m$?